CRIBS

1: (a) Show: $\tanh^{-1} x = 0.5 \ln[(1+x)/(1-x)]$ (-1<x<+1) Put y= $\tanh^{-1}x$: x= $\tanh y = (e^{y} - e^{-y})/(e^{y} + e^{-y})$ With s= e^{y} (s-1/s))/(s+1/s)=x s²=(1+x)/(1-x) Take logs: y= $\tanh^{-1} x = 0.5 \ln[(1+x)/(1-x)]$

(b) Show that $\cosh^{-1} x = \ln [(x + \sqrt{x^2 - 1})]$ $(x \ge 1)$ Put $y = \cosh^{-1} x$ $x = \cosh y = (e^y + e^{-y})/2$ With $s = e^y$ 2x = (s+1/s) $s^2 - 2sx + 1 = 0$ $s = 0.5[x \pm \sqrt{x^2 - 1}]$ $y = \cosh^{-1} x = \ln [(x + \sqrt{x^2 - 1})]$

2: $z^4 = r^4 e^{i4\theta}$ $\sqrt{z} = r^{1/2} e^{i\theta/2}$ and $\ln(z) = \ln r + i\theta$

P(i)=0 = P(-i) Divide by $(z^{2}+1)$ So $P(z)=(z^{2}+1)$ $(z^{2}+5z+6)$ The zeroes are i, -i, -2, and -3

3: Vector form of plane: $\mathbf{r}.(A,B,C)=D$. Vector normal to plane is $\mathbf{n}=(A,B,C)$

Eqn of line through (l,m,n) and perpendicular to the plane is $\mathbf{r} = (l,m,n) + t(A,B,C)$ Meets plane when $\mathbf{r}.(A,B,C) = lA + mB + nC + t(A,B,C).(A,B,C) = D$

Or: $t = [D - (lA + mB + nC)]/[A^2 + B^2 + C^2]$

The point on the plane N where this line meets the plane is $(l,m,n) + \{[D - (lA+mB+nC)]/[A^2+B^2+C^2]\}(A,B,C)$

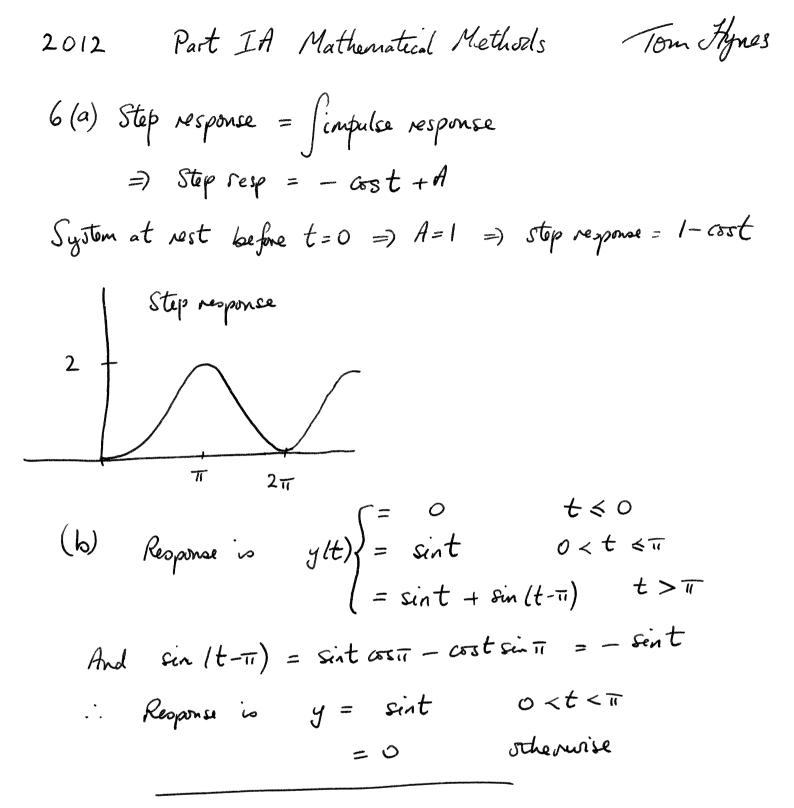
The length of PN = { $[D - (1A + mB + nC)]/[A^2 + B^2 + C^2]$ } $\sqrt{[A^2 + B^2 + C^2]}$

Particular question: (1,m,n)=(2,-1,3) (A,B,C,)=(2,-2,-1) and D=9

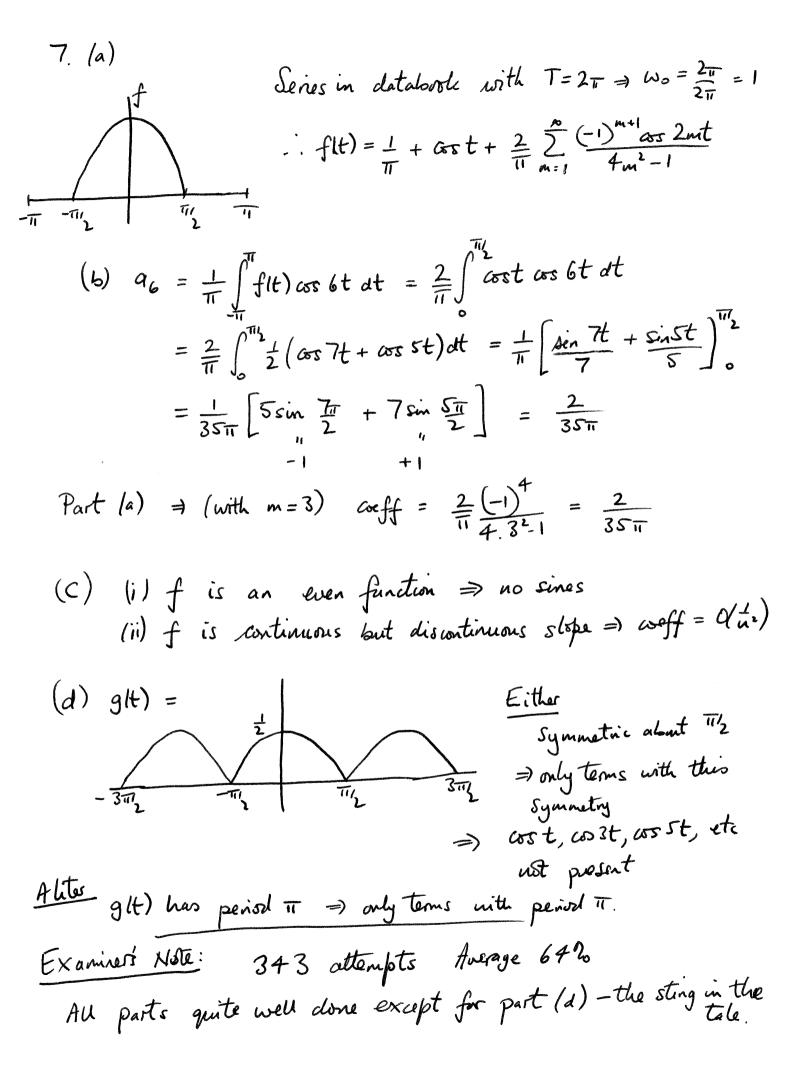
T = [9 - (4 + 2 - 3)]/9 = 2/3X 3 = 2

 $(Q_{\mathbf{z}}) = det(Q_{\mathbf{z}}) = det(\mathbf{z}) = 1$ $det Q = det Q^{0}$ => $\left(det(\underline{Q})\right)^{2} = 1 \Rightarrow det \underline{Q} = \pm 1$ => Uilij = Sij ie columns are an orthonormal set. (Also tre for rows). $(i) det (A - \Lambda I) = det \begin{pmatrix} 3/5 - \lambda & 0 & -4/5 \\ 0 & 1 - \lambda & 0 \\ -4/5 & 0 & -3/5 - \lambda \end{pmatrix} = 0$ EXPAND By 2nd row => (1-2) [(3/5-2)(-3/5-2) - 16/25] = 0 $(1-\lambda) [\lambda^2 - \frac{9}{25} - \frac{16}{25}] = 0$ $(1-\lambda)(\lambda^2-1) = 0$ $\lambda = 1, 1, -1$ Reported e-value to expect eigen-plane! $\begin{array}{c} \begin{array}{c} \begin{array}{c} \begin{array}{c} \end{array} \end{array} \end{array} \end{array} \\ \end{array} \end{array} \begin{array}{c} \begin{array}{c} \end{array} \end{array} \end{array} \\ \begin{array}{c} \end{array} \end{array} \end{array} \left(\begin{array}{c} \end{array} \end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\end{array} \right) \left(\end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\end{array} \right) \left(\end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\end{array} \right) \left(\end{array} \right) \left(\end{array} \right) \left(\begin{array}{c} \end{array} \right) \left(\bigg) \left(\bigg) \left(\\ \left) \left(\end{array} \right) \left(\bigg) \left($ (ii) A symmetric => e-values are real, e-vectors for distinct e-vals are perpendicular & repeated e-vals (e-plone) can the choose e-vectors to be perpendicular. $\begin{pmatrix} 0 \\ 1 \\ 0 \\ 1 \end{pmatrix} \begin{pmatrix} -2 \\ 0 \\ 1 \\ 0 \\ 1 \end{pmatrix} = 0 \Rightarrow \frac{1}{2} r \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \\ 2 \end{pmatrix} = 0 \bigvee \begin{pmatrix} 1 \\ 0 \\ 0 \\ 2 \end{pmatrix} = 0 \bigvee \begin{pmatrix} 1 \\ 0 \\ 0 \\ 2 \\ 1 \end{pmatrix} = 0 \bigvee$ (iii) A corresponds by reflection in a plane with normal (1,0,2) (is eventor corresponding to n=-1), other two e-values are both wit.

2=0@x=0 (200 slope) $\int_{\sqrt{1+z^2}}^{z} dz = \int_{0}^{\infty} k dx$ DATA Booh => [Sinh"(Z)] = [HX] => sinh" (do) = tex $\begin{array}{rcl} \text{REARRANSE} &=> & \frac{dy}{dx} = \sinh\left(wx\right) \\ \hline \end{array}$ $y = \frac{1}{\pi} \cosh(hx) + const.$ INTESNATE $y=0 @ x=0 \Rightarrow y= \frac{1}{k} (cosh(kx)-i) \frac{k chosen to fit-}{though(x_i, y_i)}$ y A ds_{-}/dy Pythegans $ds^{2} = dx^{2} + dy^{2}$ = dy $= ds = \sqrt{1 + (dy)^{2}}$ dx dxb) $\frac{ds}{dx} = \sqrt{1 + \sin^2(kx)} = \cosh(kx)$ $= 3 = 1 \operatorname{sinh}(kx) + \operatorname{const}_{k}$ 5=0@x=0 => 5= 1 sinh (+x) K $\frac{\text{Length}}{K} = \frac{2}{K} \frac{\sinh(\ln x_i)}{K}$ $c): |y| < < x, \Rightarrow |dy| dy < < 1 \Rightarrow \frac{d^2 y}{dx^2} = k$ (ii) $dy = kx + A \Rightarrow y = \frac{1}{2}kx^2 + Ax + B$ y=0 € x=0 & dolde=0 =) A=B=0 => y= 1/2 kx2 (iii) $y = \frac{1}{K} \left(\cosh(\frac{kx}{2}) - y \right) = \frac{1}{K} \left(\left[1 + \left(\frac{kx}{2}\right)^2 + O(x^4) \right] - 1 \right) = \frac{1}{2} \frac{kx^2}{2} + O(x^4)$ DATABOOK.



Examiner's Note 345 Candidates 344 attempted this question. Average 68%. Well done, in general. Only common error was failure to realise that the Answer to part (6) has different forms for different time intervals. Some candidates tried to do part (6) by convolution and sank without trace.



8. On one round

$$P(win) = P(2H) + P(2T)$$

 $= \frac{2}{5} \cdot \frac{3}{5} + \frac{3}{5} \cdot \frac{2}{5} = \frac{12}{25}$
 $P(loss) = 1 - P/win) = \frac{13}{25}$

For
$$\neq 1$$
 state
Expected Return = $2 \times \frac{12}{25} + 0 \times \frac{13}{25}$ (= mean = $\sum_{x} P(x=x)$)
= $\frac{24}{25}$ =) expected lars = 4p

After 10 rounds, expected loss = 40p

Examiner's Note: 335 attempts, average 66%
Average of 66% came from half the candidates producing perfect
Solutions and the other half getting nowhere. To those who
expected to losse more than £10, I would like to insite you to
Come and play!
Commonest errors:
(i) many candidates thought average of integers must be an
integer
(ii) Some thought "expected" => most probable
N.B. mean (X) is usually donted E[X]
is mean = expectation.
Two candidates earned my undying admiration by evaluating
all possible cases and conducted, after 3 pages, that
expected bits =
$$\sum_{w=0}^{\infty} {}_{10} C_w \left(\frac{12}{25}\right)^{vow} (10-2w) = 40p$$

Awesome!

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(d) From part (b)

$$y'(s) = \frac{1}{s^{\frac{1}{2}+2s^{2}}} \implies s^{3}Y + 2s^{2}Y = 1$$
Nay $L\left(\frac{d^{3}y}{dt^{3}}\right) = s^{3}Y - s^{2}y(0) - sy(0) - y(0)$

$$L\left(\frac{d^{3}y}{dt^{3}}\right) = s^{2}Y - sy(0) - y(0)s^{2} - s(y(0) + 2y(0))$$

$$= L\left(\frac{d^{3}y}{dt^{3}} + 2\frac{d^{3}y}{dt^{3}}\right) = (s^{3} + 2s^{2})Y - y(0)s^{2} - s(y(0) + 2y(0))$$

$$= \frac{y(0)}{2} + 2y(0) = 0 \qquad y(0) - 2y(0) = 0$$
Grupping this with $(s^{3} + 2s^{2})Y - 1 = 0 \Rightarrow satisfies ey?$
and $\Rightarrow y(0) = 0 \qquad y(0) + 2y(0) = 0 \qquad y(0) = 1$

$$= \frac{y(0)}{2} + \frac{y(0)}{2} = 0 \qquad y(0) = 1$$

$$= \frac{1}{2} + \frac{1}{2} +$$

(a) Rote of dange with distance in direction
$$\underline{m}$$

$$= \frac{df}{ds} = \underline{m} \cdot \nabla f \quad provided \underline{m} \quad unit water.$$
In this are
 $\nabla f = (2e^{x} - y - 1, -x - 1 + \frac{2}{(y+1)^{x}})$
At (0,0) $\nabla f = (1,1)$
 $\underline{m} = (\underline{1}, \underline{2}) \Rightarrow df = (\underline{1}, \underline{2}) \cdot (1,1) = 3$
 $\sqrt{5}$
 $\sqrt{5}$

Part IA Paper 4 Mathematical Methods

Section C

James Matheson

11 (a) There are 2^{48} or roughly 2.5×10^{14} possible MAC addresses. Even if only one byte of memory is required per table entry, this would require 2.5 TB of memory; clearly impractical! A hash table would require much less memory, typically only enough per entry (100) plus some spare to make collisions less likely, say 256 in all, i.e. a factor of 10^{12} less.

(b) We clearly need to avoid using the most significant 24 bits since these will be identical for all devices made by the same manufacturer, leading to collisions in the hash table. This leaves the least significant 24 bits of which the least significant N are most likely to be random, assuming that the manufacturer allocates addresses sequentially. If we want a hash table with 256 entries as suggested in (a) above, this implies using the least significant 8 bits. Any N less than 7 will results in unresolvable collisions and anything much larger than 9 or 10 will be a poor trade-off between wasting memory and reducing collisions, though how much this matters depends on the size of each entry in the hash table.

12 (a)
struct component {
 float cost;
 float total;
 };

(b)There are 1000 item totals, 12 monthly totals and one year total. So, as the totals accumulate, the year total will be greatest, followed by the monthly totals, and the per-item totals will generally be the smallest. As soon as any particular total gets 2^{23} times larger than the smallest quantity being added (a penny), i.e. larger than 2^{23} pence or about 80,000 pounds, individual pence will no longer be added. And once a total exceeds 2^{23} times an item value, adding the item will have no effect on the total. This is because the mantissa of the sale value has to be shifted 23 bits to the right to get equal-exponent addition. This problem is going to affect the year total first, then the monthly totals and only at most a few of item totals, which is consistent with the accountant's findings.

2012