

Engineering Tripos Part IIA, Module 3F1,
SIGNALS AND SYSTEMS
SAMPLE SOLUTIONS TO EXAM MAY 2013

1. **Solution:** A discrete-time system with input sequence $\{u_k\}$ and output sequence $\{y_k\}$, has transfer function $G(z) = \frac{(z+1)}{z(z-1)}$.

- (a) Calculate the step response of this system, and check your answer is consistent with the initial value theorem and/or the final value theorem if either one applies.

$$U(z) = z/(z-1) \text{ and}$$

$$Y(z) = G(z)U(z) = \frac{(z+1)}{(z-1)^2} = \frac{(z^{-1} + z^{-2})}{(1-z^{-1})^2} \quad \checkmark$$

and from the data book z-transform table we have $\mathcal{Z}\{k\}_{k \geq 0} = \frac{z^{-1}}{(1-z^{-1})^2}$ and $\mathcal{Z}\{k-1\}_{k \geq 1} = \frac{z^{-2}}{(1-z^{-1})^2}$ and hence $y_0 = 0$ and $y_k = 2k-1$ for $k \geq 1$. The final value theorem does not apply since the poles of $(z-1)Y(z)$ include one at $z=1$, but the initial value theorem does apply and $y_0 = \lim_{z \rightarrow \infty} Y(z) = 0$.

- (b) Write down a difference equation with this transfer function and check the first three values of the step response calculated above agree with the corresponding solution of the difference equation.

Difference equation,

$$y_{k+2} - y_{k+1} = u_{k+1} + u_k$$

has transfer function $G(z)$. With $y_k = 0$ for $k < 0$ and $u_k = 1$ for $k \geq 0$ successively solving this difference equation gives $y_0 = 0$, $y_1 = 1$ and $y_2 = y_1 + u_1 + u_0 = 3$, agreeing with the above.

- (c) Show that

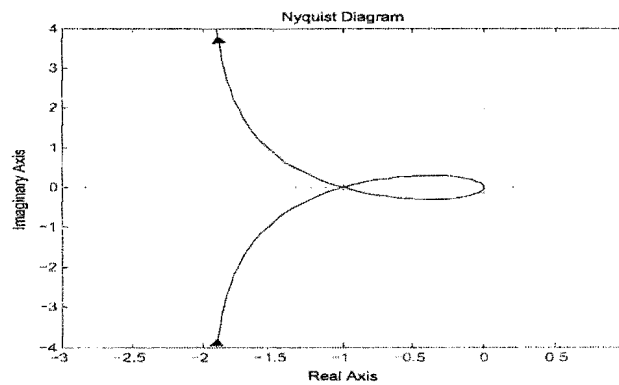
$$G(e^{j\theta}) = \frac{-j}{\tan(\theta/2)} e^{-j\theta}.$$

Note that $\frac{e^{j\theta}+1}{e^{j\theta}-1} = \frac{e^{j\theta/2}+e^{-j\theta/2}}{e^{j\theta/2}-e^{-j\theta/2}} = \frac{2\cos(\theta/2)}{2j\sin(\theta/2)}$ which immediately gives the result.

- (d) Sketch the Nyquist diagram for $G(z)$. This system is now connected in a standard unity gain negative feedback arrangement with a precompensator with a constant gain, K . Using the Nyquist stability criterion determine what values of K give closed-loop stability.

$$\begin{aligned}
G(e^{j\theta}) &= \frac{-j}{\tan(\theta/2)} e^{-j\theta} \\
&= \frac{\sin(\theta)}{\tan(\theta/2)} - \frac{j \cos(\theta)}{\tan(\theta/2)} \\
&= -2 \cos^2(\theta/2) - \frac{j \cos(\theta)}{\tan(\theta/2)} \\
&\rightarrow -2 - j\infty \text{ as } \theta \rightarrow 0
\end{aligned}$$

Note that at $G(e^{j\pi/2}) = -1$. This gives the following Nyquist plot, where the semicircle goes round to the right. There will hence be no encirclements of the point $-1/K$ if $0 < K < 1$ when the closed-loop system will be stable. There will be two clockwise encirclements if $K > 0$ and one clockwise encirclement if $K < 0$.



- (e) For $K = 1/2$ and the external reference signal, $r_k = \cos(\omega kT)$, determine the behaviour of the error, $e_k = r_k - y_k$ as k becomes large for the three cases, $\omega T = 0, \pi/4, \pi$.

$EX(z) = \frac{1}{1+KG(z)}R(z)$ and since $K = 1/2$ is stabilising we have that $y_k \rightarrow |H(e^{j\theta})| \cos(\omega kT + \angle G(e^{j\theta}))$, where $H(z) = \frac{1}{1+KG(z)}$, $\theta = \omega T = 0, \pi/4, \pi$.
Now

- (i) since $G(e^{j0}) = \infty$, $H(e^{j0}) = 0$, so $e_k \rightarrow 0$;
- (ii) $G(e^{j\pi}) = 0$, $H(e^{j\pi}) = 1$, so $e_k \rightarrow (-1)^k$;
- (iii) $G(e^{j\pi/4}) = \frac{-j \cos(\pi/4)}{\tan(\pi/8)} - \frac{\sin(\pi/4)}{\tan(\pi/8)} = -1.71 - 1.71j$, $H(e^{j\pi/4}) = \frac{1}{0.146 - 0.854j} = 1.15e^{1.40j}$, so $e_k \rightarrow 1.15 \cos(\omega kT + 1.40)$

(f) If the input to $G(z)$ is $u_k = (-1)^k$ for $k \geq 0$, calculate the output and comment on its relation to the frequency response.

In this case $U(z) = \frac{z}{z+1}$ so $Y(z) = \frac{1}{z-1}$ and hence $y_0 = 0$ and $y_k = 1$ for $k \geq 1$. Hence the output due to the sinusoidal input is zero because $G(e^{j\pi}) = 0$ but the output does not tend to zero because the system is not stable due to the pole at $z = 1$.

2. Solution:

(a) A linear discrete-time system with input sequence $\{u_k\}$ and output sequence $\{y_k\}$, has pulse response sequence $\{g_k\}$ and transfer function $G(z)$.

i. Show that if

$$\sum_{k=0}^{\infty} |g_k| = M < \infty \quad (1)$$

then bounded inputs will produce bounded outputs.

Standard bookwork.

ii. In the case $G(z) = \frac{1}{(z^2 + 1)}$ show that (1) does not hold and that there exists a bounded input that gives an unbounded output sequence.

For $G(z) = \frac{1}{z^2+1}$, the pulse response will be $g_0 = 0, g_1 = 0, g_2 = 1, g_3 = 0, g_4 = -1, g_5 = 0, g_6 = 1, \dots$. Or from the databook with $\omega_0 T = \pi/2$, $r = 1$, $a = 0$, giving $g_k = \sin(\pi(k-1)/2)$, for $k \geq 2$. Clearly, although the g_k remain bounded the $\sum |g_k| \rightarrow \infty$. If the input sequence, $u_k = g_k$ then $Y(z) = \frac{1}{(z^2+1)^2} = z^{-2}(1+z^{-2})^{-2} = z^{-2}(1-2z^{-2}+3z^{-4}-4z^{-6}+\dots(-1)^k(k+1)z^{-2k}+\dots)$. This clearly increases without bound as $k \rightarrow \infty$, giving an unbounded output for a bounded input.

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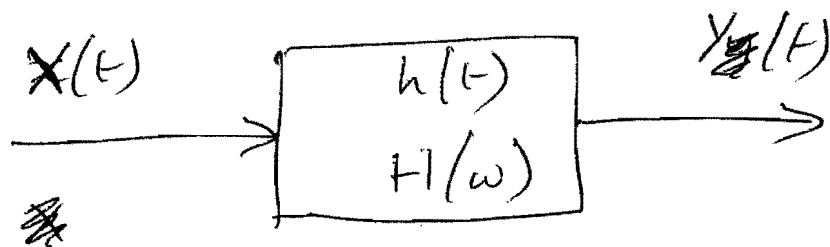
- (c) Show that

$$G(e^{j\theta}) = \frac{-j}{\tan(\theta/2)} e^{-j\theta}.$$

Note that $\frac{e^{j\theta} + 1}{e^{j\theta} - 1} = \frac{e^{j\theta/2} + e^{-j\theta/2}}{e^{j\theta/2} - e^{-j\theta/2}} = \frac{2 \cos(\theta/2)}{2j \sin(\theta/2)}$ which immediately gives the result.

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2. (b)



$$(i) \tau_{xx}(z) = E[x(t) x(t+z)]$$

$$y(t) = x(t) * h(t) \quad (\text{convolution})$$

$$= \int_{-\infty}^{\infty} h(\beta) x(t-\beta) d\beta$$

$$\tau_{xy}(z) = E[x(t) y(t+z)]$$

$$= E\left[x(t) \int_{-\infty}^{\infty} h(\beta) x(t+z-\beta) d\beta\right]$$

$$= \int_{-\infty}^{\infty} h(\beta) E[x(t) x(t+z-\beta)] d\beta$$

$$= \int_{-\infty}^{\infty} h(\beta) \tau_{xx}(z-\beta) d\beta$$

$$= \underline{\underline{h(z) * \tau_{xx}(z)}} \quad (\text{convolution})$$

2 b. (ii) $S_x(\omega)$ is the Fourier transform of $r_{xx}(\tau)$.

$S_{xy}(\omega)$ is the Fourier transform of $r_{xy}(\tau)$

(iii) Taking the Fourier transform of the convolution result in b(i):

$$S_{xy}(\omega) = H(\omega) \cdot S_{xx}(\omega)$$

$$\therefore H(\omega) = \frac{S_{xy}(\omega)}{S_{xx}(\omega)} = \frac{\text{F.T.} \{ r_{xy}(\tau) \}}{\text{F.T.} \{ r_{xx}(\tau) \}}$$

In this way $H(\omega)$ may be calculated when the input to the system comprises random perturbations which typically occur in normal operation (eg an airliner cruising at altitude & subject to atmospheric turbulence, or a power station subject to random load fluctuations)

$$3. (a) \quad \Phi_x(u) = E[e^{jux}]$$

$$= \int_{-\infty}^{\infty} e^{jux} f_x(x) dx$$

whereas the Fourier transform of $f_x(t)$ is

$$F_x(\omega) = \int_{-\infty}^{\infty} f_x(t) e^{-j\omega t} dt$$

Hence exchanging x for t , & u for $(-\omega)$, the two formulae are equivalent.

So the characteristic function $\Phi_x(u) = F_x(-u)$ where $F_x(\cdot)$ is the Fourier transform of $f_x(x)$, the pdf.

$$(b) \quad Y = X_1 + X_2$$

Now joint pdf $f(y, x_1) = f(y/x_1) \cdot f_1(x_1)$

$$\text{Now } f(y/x_1) = f_2(y - x_1)$$

∴ marginalising out x_1 , gives

$$f(y) = \int_{-\infty}^{\infty} f(y/x_1) f_1(x_1) dx_1,$$

$$= \int_{-\infty}^{\infty} f_2(y - x_1) f_1(x_1) dx_1,$$

$$= f_2 * f_1 \quad (\text{convolution})$$

3 (b) - cont.)

Taking Fourier transforms

$$F_y(\omega) = F_{x_2}(\omega) \cdot F_{x_1}(\omega)$$

∴ replacing ω with $-\omega$.

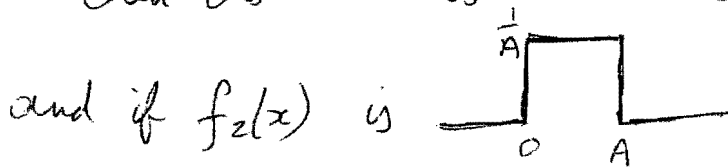
$$\Phi_x(\omega) = \Phi_{x_2}(\omega) \cdot \Phi_{x_1}(\omega)$$

Hence adding x_1 & x_2 results in multiplication of their characteristic functions.

(c) If $f_1(x)$ is



then its F.T. is $\text{sinc}^2\left(\frac{\omega A}{2}\right)$ (from Data Book)



then its F.T. is $\text{sinc}\left(\frac{\omega A}{2}\right) \cdot e^{-j\omega A/2}$

Letting $\omega = -\omega$:

$$\therefore \Phi_{x_1}(\omega) = \text{sinc}^2\left(\frac{-\omega A}{2}\right) = \text{sinc}^2\left(\frac{\omega A}{2}\right)$$

$$\text{∴ } \Phi_{x_2}(\omega) = \text{sinc}\left(\frac{-\omega A}{2}\right) \cdot e^{j\omega A/2} = \text{sinc}\left(\frac{\omega A}{2}\right) \cdot e^{j\omega A/2}$$

$$\therefore \Phi_y(\omega) = \underline{\underline{\text{sinc}^3\left(\frac{\omega A}{2}\right) \cdot e^{j\omega A/2}}}$$

$$3. (d) \quad \text{Since } \Phi_Y(u) = \int_{-\infty}^{\infty} f_Y(y) e^{juy} dy$$

$$\frac{d}{du} \Phi_Y(u) = \int_{-\infty}^{\infty} jy \cdot f_Y(y) e^{juy} dy$$

$$\frac{d^2}{du^2} \Phi_Y = \int_{-\infty}^{\infty} (jy)^2 f_Y(y) e^{juy} dy$$

$$\vdots$$

$$\frac{d^n}{du^n} \Phi_Y = \int_{-\infty}^{\infty} (jy)^n f_Y(y) e^{juy} dy$$

At the point $u=0$:

$$\left. \frac{d^n}{du^n} \Phi_Y \right|_{u=0} = j^n \int_{-\infty}^{\infty} y^n f_Y(y) dy$$

since $e^{j0 \cdot y} = 1$
at $u=0$.

$$= j^n \cdot (n^{\text{th}} \text{-order moment of } f_Y(y))$$

For the above case, with $n=1$,

~~$$\frac{d}{du} \Phi_Y(u) = \frac{d}{du} \left[\frac{\sin^3(uA/2)}{(uA/2)^3} \cdot e^{j u A/2} \right]$$~~

3 (d - cont)

For the above case with $n=1$

$$\begin{aligned}\frac{d}{du} \Phi_Y &= \frac{d}{du} \left[\text{sinc}^3\left(\frac{uA}{2}\right) \cdot e^{juA/2} \right] \\ &= 3 \text{sinc}^2\left(\frac{uA}{2}\right) \cdot \frac{d}{du} \left(\text{sinc}\left(\frac{uA}{2}\right) \right) \cdot e^{juA/2} \\ &\quad + \text{sinc}^3\left(\frac{uA}{2}\right) \cdot \cancel{j} \frac{jA}{2} \cdot e^{juA/2}\end{aligned}$$

Now at $u=0$, $\frac{d}{du} \left(\text{sinc}\left(\frac{uA}{2}\right) \right) = 0$ since the sinc function has zero gradient at its centre point, and $\text{sinc}\left(\frac{uA}{2}\right) = 1$ at this point.

$$\therefore \left. \frac{d}{du} \Phi_Y \right|_{u=0} = 0 + (1)^3 \cdot \frac{jA}{2} \cdot 1 = \frac{jA}{2}$$

Hence the first-order moment (or mean value) of $f_Y(y)$ is $\frac{A}{2}$, when we divide $\frac{d}{du} \Phi_Y$ by j .

This is what we would expect as the mean of X_1 is zero & the mean of X_2 is $\frac{A}{2}$, by inspection of the pdfs, f_1 & f_2 , and $Y = X_1 + X_2$.

The matrix for a

4. (a) ~~A~~ valid joint pdf should sum to unity.

The columns of $P(X_{n-1}, X_n)$ sum to 0.3, 0.4, 0.3 & hence the matrix does sum to unity, and is valid.

The matrix is symmetric, so both the rows & ~~the~~ columns sum to $[0.3 \ 0.4 \ 0.3]$ & hence this gives the probabilities for the 3 states, A, B and C, respectively.

$$\therefore P(X_n) = \left. \begin{array}{l} P(X_n=A) = 0.3 \\ P(X_n=B) = 0.4 \\ P(X_n=C) = 0.3 \end{array} \right\} \text{ for any } n.$$

(b) The conditional table defines $P(X_n | X_{n-1})$

$$\# \text{ But } P(X_n, X_{n-1}) = P(X_n | X_{n-1}) \cdot P(X_{n-1})$$

$$\therefore P(X_n | X_{n-1}=A) = \frac{P(X_n, X_{n-1}=A)}{P(X_{n-1}=A)} = \begin{bmatrix} 0.2 \\ 0.05 \\ 0.05 \end{bmatrix} \cdot \frac{1}{0.3}$$

Similarly for $X_{n-1}=B$ & C:

$$\therefore P(X_n | X_{n-1}) = \begin{bmatrix} 0.667 & 0.125 & 0.167 \\ 0.167 & 0.75 & 0.167 \\ 0.167 & 0.125 & 0.667 \end{bmatrix} = \begin{bmatrix} \frac{2}{3} & \frac{1}{8} & \frac{1}{6} \\ \frac{1}{6} & \frac{3}{4} & \frac{1}{6} \\ \frac{1}{6} & \frac{1}{8} & \frac{2}{3} \end{bmatrix} \quad \left(\begin{array}{l} \text{fractions} \\ \text{are} \\ \text{more} \\ \text{accurate} \end{array} \right)$$

4. (b) (cont.)

Mutual info: $I(X_n; X_{n-1}) = H(X_n) - H(X_n | X_{n-1})$

$$H(X_n) = - \sum_{i=1}^3 p_i \log_2 p_i = - (2 \cdot 0.3 \log_2 0.3 + 0.4 \log_2 0.4) \\ = 2 \cdot 0.5211 + 0.5288 = \underline{1.5710} \text{ bit/symbol}$$

$$H(X_n | X_{n-1} = A) = - \left(2 \cdot \frac{1}{6} \log_2 \frac{1}{6} + \frac{2}{3} \log_2 \frac{2}{3} \right) \\ = 2 \cdot 0.4308 + 0.3900 = 1.2516 \text{ bit/symbol}$$

$$H(X_n | X_{n-1} = B) = - \left(2 \cdot \frac{1}{8} \log_2 \frac{1}{8} + \frac{3}{4} \log_2 \frac{3}{4} \right) \\ = 2 \cdot 0.3750 + 0.3113 = 1.0613 \text{ bit/symbol}$$

$$H(X_n | X_{n-1} = C) = H(X_n | X_{n-1} = A) = 1.2516 \text{ bit/symbol}$$

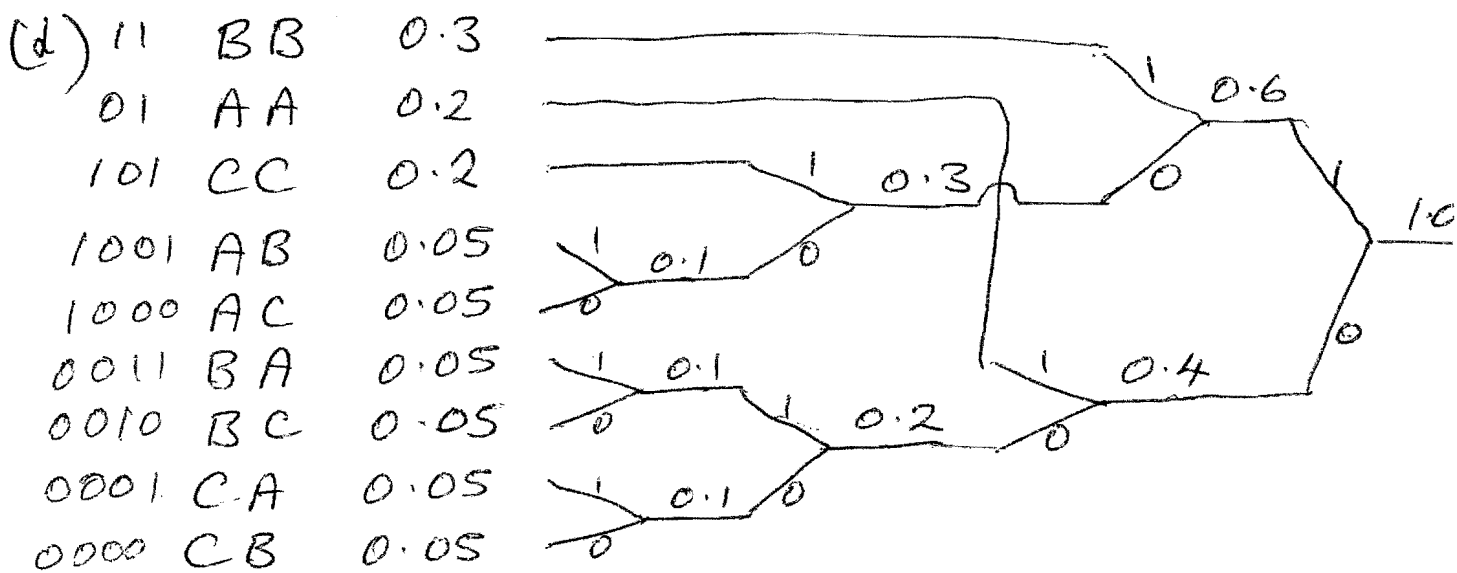
$$\therefore H(X_n | X_{n-1}) = 2 \cdot 0.3 \cdot 1.2516 + 0.4 \cdot 1.0613 \\ = 2 \cdot 0.3755 + 0.4245 = \underline{1.1755} \text{ bit/sym}$$

$$\therefore I(X_n; X_{n-1}) = 1.5710 - 1.1755 = \underline{\underline{0.3955}} \\ \text{bit/sym}$$

4. (c) For a block of 4 symbols:

$$\begin{aligned}
 \text{Total entropy} &= H(X_{n-3}) + H(X_{n-2}|X_{n-3}) \\
 &\quad + H(X_{n-1}|X_{n-2}) + H(X_n|X_{n-1}) \\
 &= H(X_n) + 3H(X_n|X_{n-1}) \\
 &= 4H(X_n) - 3I(X_n; X_{n-1}) \\
 &= 4 \cdot 1.5710 - 3 \cdot 0.3955 = 6.2840 - 1.1865 \\
 &= \underline{5.0975} \text{ bits/block.}
 \end{aligned}$$

This is the average no. of bits to encode each block, assuming an ideal encoder.



$$\begin{aligned}
 \text{Av. no of bits to code 2 symbols} &= 2(0.3+0.2) + 3 \cdot 0.2 \\
 &\quad + 4 \cdot 6 \cdot 0.05 \\
 &= 1.0 + 0.6 + 1.2 = 2.8 \text{ bit/(2 sym)}
 \end{aligned}$$

$$\therefore \text{Av no of bits to code 4 symbols} = 2 \cdot 2.8 = \underline{5.6} \text{ bits}$$

$$\text{Efficiency} = \frac{\text{Entropy}}{\text{Av no. of bits}} = \frac{5.0975}{5.6} = \underline{91.03\%}$$

t. e) The main loss of efficiency in (d) is because we have not taken account of the mutual info between ~~the~~ ^{3 symbols} 2 & 3 of the 4-bit block. This would reduce the word length by approx 0.3955 bits to $5.6 - 0.3955 = 5.2045$ bits/block, which would increase the efficiency to $\frac{5.0975}{5.2045} = \underline{\underline{97.9\%}}$

Arithmetic coding would probably be the easiest way to increase the efficiency in practice, since one could start each ~~the~~ block with $P(X_n)_{\text{coding}} = [0.3 \ 0.4 \ 0.3]^T$ as the prob distribution for the first symbol, and then use appropriate columns from $P(X_n | X_{n-1})$ for each subsequent symbol X_n , given knowledge of the previous symbol X_{n-1} .

This should get a result very close to the calculated entropy of 5.0975 bit/block = 1.2744 bit/sym.

If one used Arithmetic coding over much longer blocks, then one should be able to approach closely the conditional entropy, $H(X_n | X_{n-1}) = \underline{\underline{1.1755}}$ bit/sym.